TYPES 2019

Weak Type Theory

a rather strong type theory

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Conversion

Extends the notion of β-equality

$$(\lambda x.t)$$
 $u \equiv t[x \leftarrow u]$

Conversion

Extends the notion of β -equality

$$(\lambda x.t)$$
 $u = t[x \leftarrow u]$

$$vec_N$$
 (3 + 2) $\equiv vec_N$ 5

Conversion

Extends the notion of β -equality

$$(\lambda x.t)$$
 $u \equiv t[x \leftarrow u]$

$$vec_{\mathbb{N}} (3 + 2) \equiv vec_{\mathbb{N}} 5$$
[1; 2; 3] ++ [4; 5] \equiv [1; 2; 3; 4; 5]

Witness property

```
from \exists x, Px
compute X
```

Witness property

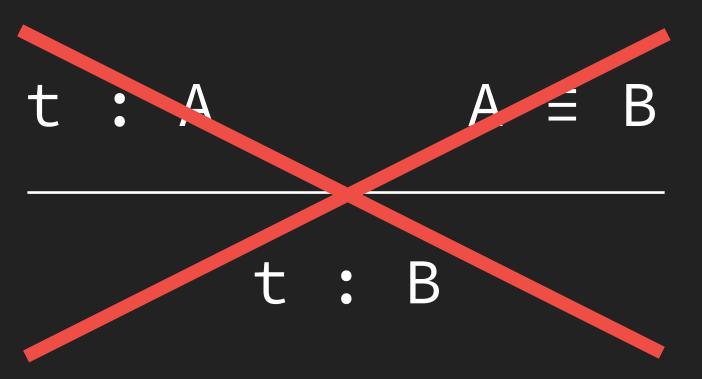
```
from \exists x, Px
compute X
```

```
(t: \exists x, Px) \rightarrow \langle u; h \rangle
```

Witness property

```
from \exists x, Px
compute X
```

```
(t: \exists x, Px) \rightarrow \langle u; h \rangle
with h: Pu
```



$$u =_A v$$

$$u =_A v$$

$$refl_A x : x =_A x$$

```
u =_{A} V
refl_{A} x : x =_{A} x
coe_{A} : \forall x y, x =_{A} y \Rightarrow \forall P, P x \Rightarrow P y
uip_{A} : \forall p q, p =_{x = y} q
```

```
u =_{A} V
refl_{A} x : x =_{A} x
coe_{A} : \forall x y, x =_{A} y \Rightarrow \forall P, P x \Rightarrow P y
uip_{A} : \forall p q, p =_{x = y} q
\beta t u : (\lambda x.t) u = t[x := u]
```

Theory without computation

```
u =_{A} V
refl_{A} x : x =_{A} x
coe_{A} : \forall x y, x =_{A} y \Rightarrow \forall P, P x \Rightarrow P y
uip_{A} : \forall p q, p =_{x = y} q
\beta t u : (\lambda x.t) u = t[x := u]
```

and more...

Congruence for free

```
coe<sub>A</sub>: \forall x y, x =_A y \rightarrow \forall P, P x \rightarrow P
```

Congruence for free

coe_A:
$$\forall x y, x =_A y \rightarrow \forall P, P x \rightarrow P y$$

$$x = y \rightarrow y = z \rightarrow x = z$$

Congruence for free

coe_A:
$$\forall x y, x =_A y \Rightarrow \forall P, P x \Rightarrow P y$$

$$x = y \Rightarrow y = z \Rightarrow x = z$$

$$f = g \Rightarrow x = y \Rightarrow f x = g y$$

Congruence for free ??

coe_A: $\forall x y, x =_A y \rightarrow \forall P, P x \rightarrow P y$

Congruence for free ??

coe_A:
$$\forall x y, x =_A y \rightarrow \forall P, P x \rightarrow P y$$

Problems under binders

$$\vdash A = A'$$
 $\times : A \vdash B = B'$

$$\vdash \Pi(x:A).B = \Pi(x:A').B'$$

Congruence for free ??

coe_A:
$$\forall x y, x =_A y \rightarrow \forall P, P x \rightarrow P y$$

Problems under binders

$$\vdash A = A'$$
 $\times : A \vdash B = B'$

$$\vdash \Pi(x:A).B = \Pi(x:A').B'$$

⇒ more axioms needed

ETT ITT

ETT WTT WTT

no computation

no conversion

ETT

ITT

WTT

$$p : u = v$$

 $u \equiv v$

no computation

no conversion

equality reflection

ETT ITT DIE WITT

equality reflection

 $u \equiv v$

no conversion

ETT WTT WTT

ETT WITH WITH

terms up to computation

ETT ITT WTT

terms up to computation

ETT ITT WTT

up to equality terms up to computation





up to equality terms up to computation decidable type checking trivial checking

